

MTL Algebra of Fractions and Maximal *MTL* Algebra of Quotients

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Received: April 8, 2013 Accepted: May 16, 2013 Online Published: May 28, 2013

doi:10.5539/jmr.v5n2p115 URL: <http://dx.doi.org/10.5539/jmr.v5n2p115>

Abstract

In this paper we introduce the notions of *MTL algebra of fractions* and *maximal MTL algebra of quotients for a MTL algebra* and prove constructively the existence of a maximal *MTL algebra of quotients* (see Buşneag & Piciu, 2005, for *BL* algebras).

Keywords: residuated lattice, Boolean algebra, *MTL* algebra, *BL* algebra, multiplier, *MTL algebra of fractions*, maximal *MTL algebra of quotients*

AMS Subject Classification 2010: 06D72, 03G25

1. Introduction

A *localization ring* $A_{\mathcal{F}}$ associated with a Gabriel topology $\mathcal{F}$ for a ring A is a very important construction in ring theory. For the term *localization* we have in view Chapter IV: *Localization* in N. Popescu's book (1971). The notion of *complete ring of quotients* for a commutative ring is introduced in Lambek's book (1966). This localization is relative to the *dense ideals* and is a special case of localization ring. Schmid define in 1980, the concept of *maximal lattice of quotients* for a distributive lattice using partial morphisms introduced by Findlay and Lambek (1966). The *multipliers* (defined for a distributive lattice by W. H. Cornish in 1974 and 1980) plays an important role in this constructions.

Basic (Fuzzy) logic (*BL* from now on) is the many-valued residuated logic introduced by Hájek in 1998 to cope with the logic of continuous t-norms and their residua. *Monoidal logic* (*ML* from now on), introduced by Höhle (1995), is a logic whose algebraic counterpart is the class of residuated lattices; *MTL* algebras (see Esteva & Godo, 2001) are algebraic structures for the Esteva-Godo monoidal t-norm based logic (*MTL*), a many-valued propositional calculus that formalizes the structure of the real unit interval $[0, 1]$, induced by a left-continuous t-norm. *MTL* algebras were independently introduced in Flondor, Georgescu, and Iorgulescu (2001) under the name *weak-BL algebras*. The results obtained in this paper for *MTL* algebras are analogously to the ones obtained for *BL* algebras in Buşneag and Piciu (2005). The main difference is that the equation $x \odot (x \rightarrow y) = x \wedge y$ is not valid for *MTL* algebras.

This paper is organized as follows: Section 2 is dedicated to basic definitions and rules of calculus in *MTL* algebras. In Section 3 we introduce the notion of *multiplier* for a *MTL* algebra. In the proof of Lemma 9 and Lemma 10 we have used mainly the rules c_{13} and c_{16} which are specific for *MTL* algebras (by Proposition 4 and Corollary 5). This explain why in this paper we have considered the particular case of *MTL* algebras and not the general case of residuated lattice.

In Section 4 we introduce the notions of *MTL algebra of fractions* and *maximal MTL algebra of quotients for a MTL algebra*. In Theorem 30 we prove the existence of the maximal *MTL algebra of quotients for a MTL algebra*.

This paper is a very important step in a future study of localization in the category of *MTL* algebras (and more general, in the category of residuated lattices).

For a survey relative to notion of fractions and localization in algebra of logic see Rudeanu (2010).

2. Definitions and First Properties

In this section we review the basic results relative to *MTL* algebras with more details and examples.

Definition 1 An algebra $(L, \wedge, \vee, \odot, \rightarrow, 0, 1)$ of type $(2, 2, 2, 2, 0, 0)$ equipped with an order $\leq$ is a residuated lattice (Blyth & Janovitz, 1972; Galatos, Jipsen, Kowalski, & Ono, 2007; Turunen, 1999), if it satisfies:

- (LR_1) $(L, \wedge, \vee, 0, 1)$ is a bounded lattice relative to order $\leq$;
- (LR_2) $(L, \odot, 1)$ is an ordered commutative monoid;
- (LR_3) $(\odot, \rightarrow)$ is an adjoint pair ($z \leq x \rightarrow y$ iff $x \odot z \leq y$ for every $x, y, z \in L$).

For examples of residuated lattices see Buşneag and Piciu (2006), Galatos et al. (2007), and Turunen (1999).

In this section by L we denote the universe of a residuated lattice. We denote $x^* = x \rightarrow 0$ and $x^{**} = (x^*)^*$, for $x \in L$.

We review some rules of calculus for residuated lattices:

Theorem 1 (Buşneag & Piciu, 2006; Galatos et al., 2007) *Let $x, y, z \in L$. Then:*

- (c_1) $x \rightarrow x = 1, 1 \rightarrow x = x, 0 \rightarrow x = 1, y \leq x \rightarrow y, x \odot (x \rightarrow y) \leq y, x \rightarrow 1 = 1, x \odot 0 = 0$;
- (c_2) $x \leq y$ iff $x \rightarrow y = 1$;
- (c_3) $x \leq y$ implies $x \odot z \leq y \odot z, z \rightarrow x \leq z \rightarrow y$ and $y \rightarrow z \leq x \rightarrow z$;
- (c_4) $x \rightarrow (y \rightarrow z) = (x \odot y) \rightarrow z = y \rightarrow (x \rightarrow z)$, so $(x \odot y)^* = x \rightarrow y^* = y \rightarrow x^*$;
- (c_5) $x \odot x^* = 0$ and $x \odot y = 0$ iff $x \leq y^*$;
- (c_6) $x \odot (y \vee z) = (x \odot y) \vee (x \odot z)$;
- (c_7) $x \rightarrow (y \wedge z) = (x \rightarrow y) \wedge (x \rightarrow z)$.

We shall denote $B(L) = \{x \in L : x \text{ is a complemented element in } (L, \wedge, \vee, 0, 1)\}$, which is a Boolean algebra (called the *Boolean center* of L).

Theorem 2 (Buşneag & Piciu, 2006) *For $a \in L, a \in B(L)$ iff $a \vee a^* = 1$.*

Theorem 3 (Buşneag & Piciu, 2006; Galatos et al., 2007) *If $a_1, a_2 \in B(L)$ and $x, y \in L$, then:*

- (c_8) $a_1 \odot x = a_1 \wedge x$;
- (c_9) $x \odot (x \rightarrow a_1) = a_1 \wedge x, a_1 \odot (a_1 \rightarrow x) = a_1 \wedge x$;
- (c_{10}) $a_1 \odot (x \rightarrow y) = a_1 \odot [(a_1 \odot x) \rightarrow (a_1 \odot y)]$;
- (c_{11}) $x \odot (a_1 \rightarrow a_2) = x \odot [(x \odot a_1) \rightarrow (x \odot a_2)]$.

Definition 2 (Esteva & Godo, 2001) A *MTL* algebra is a residuated lattice satisfying the prelinarity equation:

- (c_{12}) $(x \rightarrow y) \vee (y \rightarrow x) = 1$.

The variety of *MTL* algebras will be denoted by $\mathcal{MTL}$.

Example 1 (Iorgulescu, 2004) Let $L = \{0, a, b, c, d, 1\}$, with $0 < a, b < c < 1, 0 < b < d < 1$, but a, b and, respective c, d are incomparable. Then $(L, \wedge, \vee, \odot, \rightarrow, 0, 1)$ is an *MTL* algebra, where the operations $\odot$ and $\rightarrow$ are defined as follows:

$\rightarrow$	0	a	b	c	d	1
0	1	1	1	1	1	1
a	d	1	d	1	d	1
b	c	c	1	1	1	1
c	b	c	d	1	d	1
d	a	a	c	c	1	1
1	0	a	b	c	d	1

$\odot$	0	a	b	c	d	1
0	0	0	0	0	0	0
a	0	a	0	a	0	a
b	0	0	0	0	b	b
c	0	a	0	a	b	c
d	0	0	b	b	d	d
1	0	a	b	c	d	1

Proposition 4 (Esteva & Godo, 2001) *Let L be a residuated lattice. The following conditions are equivalent:*

- (i) $L \in \mathcal{MTL}$;

(ii) L is a subdirect product of linearly ordered residuated lattices;

(iii) $(c_{13}) x \rightarrow (y \vee z) = (x \rightarrow y) \vee (x \rightarrow z)$, for any $x, y, z \in L$;

(iv) $(c_{14}) (x \wedge y) \rightarrow z = (x \rightarrow z) \vee (y \rightarrow z)$, for any $x, y, z \in L$.

Corollary 5 (Esteva & Godo, 2001; Flondor, Georgescu, & Iorgulescu, 2001) Let $L \in \mathcal{MTL}$. For every $x, y, z \in L$:

$$(c_{15}) (x \wedge y)^* = x^* \vee y^*;$$

$$(c_{16}) x \odot (y \wedge z) = (x \odot y) \wedge (x \odot z);$$

$$(c_{17}) x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z);$$

$$(c_{18}) x \vee y = ((x \rightarrow y) \rightarrow y) \wedge ((y \rightarrow x) \rightarrow x).$$

Remark 1 A MTL algebra L is a BL algebra iff in L is verified the divisibility condition: $x \odot (x \rightarrow y) = x \wedge y$. So, BL algebras are examples of MTL algebras; for an example of MTL algebra which is not BL algebra see Turunen (1999, p. 16). Also, every linearly ordered residuated lattice is a MTL algebra.

3. Multipliers on a MTL Algebra

By L we denote the universe of a MTL algebra.

Let $\mathcal{I}_d(L) = \{I : I \text{ is an ideal in the lattice } (L, \wedge, \vee, 0, 1)\}$ (see Balbes & Dwinger, 1974) and $\mathcal{I}(L)$ the set of all decreasing subsets of L . We have that, $\mathcal{I}(L) \subseteq \mathcal{I}_d(L)$ and if $J_1, J_2 \in \mathcal{I}(L)$, then $J_1 \cap J_2 \in \mathcal{I}(L)$. Also, if $J \in \mathcal{I}(L)$, then $0 \in J$.

Definition 3 A map $p : J \rightarrow L$, with $J \in \mathcal{I}(L)$, is a partial multiplier on L if it verifies the axioms:

$$(M_1) p(a \odot x) = a \odot p(x), a \in B(L), x \in J;$$

$$(M_2) x \odot (x \rightarrow p(x)) = p(x), x \in J;$$

$$(M_3) \text{ If } a \in B(L) \cap J, \text{ then } p(a) \in B(L);$$

$$(M_4) x \wedge p(a) = a \wedge p(x), a \in B(L) \cap J, x \in J.$$

Remark 2 Since $x \odot (x \rightarrow p(x)) \leq x$, from (M_2) we conclude that $p(x) \leq x$, for $x \in J$.

Remark 3 We use *multiplier* instead *partial multiplier*.

By $d(p) \in \mathcal{I}(L)$ we denote the domain of p ; we call p total if $d(p) = L$.

Example 2 Let $a \in B(L)$ and $J \in \mathcal{I}(L)$. Then the map $p_a : J \rightarrow L$, $p_a(x) = a \wedge x \stackrel{(c_8)}{=} a \odot x$, for every $x \in J$ is a multiplier on L . We called this multiplier *principal*.

The axioms (M_1) , (M_3) and (M_4) are verified as in the case of BL algebras (see Buşneag & Piciu, 2005). Also, for $x \in J$, $x \odot (x \rightarrow p_a(x)) = x \odot (x \rightarrow (a \wedge x)) \stackrel{(c_7)}{=} x \odot [(x \rightarrow a) \wedge (x \rightarrow x)] = x \odot (x \rightarrow a) \stackrel{(c_9)}{=} a \wedge x = p_a(x)$, hence (M_2) is verified.

We denote p_a by $\overline{p_a}$ if $d(p_a) = L$. In particular, for $a = 0, 1$ the maps $\overline{p_0} = \mathbf{0} : L \rightarrow L$, $\overline{p_0}(x) = \mathbf{0}(x) = 0$, for every $x \in L$ and $\overline{p_1} = \mathbf{1} : L \rightarrow L$, $\overline{p_1}(x) = \mathbf{1}(x) = x$, for every $x \in L$ are total multipliers on L .

Remark 4 From (M_4) , if $J = L$, then for $a = 1$ we deduce that $x \wedge p(1) = p(x)$, so every total multiplier is principal.

For $a \in L$ and $J = [a] = \{x \in L : x \leq a\} \in \mathcal{I}(L)$ we consider the map $g_a : J \rightarrow L$, $g_a(x) = a \odot (a \rightarrow x)$ for every $x \in J$.

Lemma 6 g_a verify (M_1) , (M_3) and (M_4) .

Proof. (M_1) . For $x \in J$ and $e \in B(L) \cap J$ (hence $x \leq a, e \in B(L)$ and $e \leq a$) we have: $g_a(e \odot x) = a \odot (a \rightarrow (e \odot x)) \stackrel{(c_8)}{=} a \odot (a \rightarrow (e \wedge x)) \stackrel{(c_7)}{=} a \odot [(a \rightarrow e) \wedge (a \rightarrow x)] \stackrel{(c_{16})}{=} [a \odot (a \rightarrow e)] \wedge [a \odot (a \rightarrow x)] = (a \wedge e) \wedge g_a(x) = e \wedge g_a(x) = e \odot g_a(x)$.

(M_3) . If $e \in B(L) \cap J$, then $e \in B(L)$ and $e \leq a$, hence $g_a(e) = a \odot (a \rightarrow e) = a \wedge e = e \in B(L)$.

(M_4) . Consider $x \in J$ and $e \in B(L) \cap J$ (that is, $x, e \leq a$ and $e \in B(L)$). Thus, $e \wedge g_a(x) = e \wedge [a \odot (a \rightarrow x)] = e \odot a \odot (a \rightarrow x) = (e \wedge a) \odot (a \rightarrow x) = e \odot (a \rightarrow x)$ and $x \wedge g_a(e) = x \wedge [a \odot (a \rightarrow e)] = x \wedge (a \wedge e) = x \wedge e = e \odot x$. Since $x \leq a \rightarrow x$, then $e \odot x \leq e \odot (a \rightarrow x)$, hence $x \wedge g_a(e) \leq e \wedge g_a(x)$.

From $e \leq a$ we deduce that $a \rightarrow x \leq e \rightarrow x$ hence $e \odot (a \rightarrow x) \leq x$. Then $e \odot (a \rightarrow x) \leq e \wedge x = e \odot x$, hence $e \wedge$

$$g_a(x) = x \wedge g_a(e).$$

□

Following Lemma 6, we can obtain an example of multiplier which is not principal.

For this, we consider $L = [0, 1]$ (see Turunen, 1999, p. 16) and for all $x, y \in L$ we define

$$x \odot y = 0, \text{ if } x + y \leq \frac{1}{2} \text{ and } x \odot y = x \wedge y, \text{ if } x + y > \frac{1}{2}$$

$$x \rightarrow y = 1 \text{ if } x \leq y \text{ and } x \rightarrow y = \max\{\frac{1}{2} - x, y\} \text{ if } x > y.$$

Then $(L, \wedge, \vee, \odot, \rightarrow, 0, 1)$ is a MTL -algebra. Obviously, L is not a BL -algebra and $B(L) = \{0, 1\}$.

Lemma 7 $g_{\frac{1}{3}}: (\frac{1}{3}) = [0, \frac{1}{3}] \rightarrow L = [0, 1], g_{\frac{1}{3}}(x) = \frac{1}{3} \odot (\frac{1}{3} \rightarrow x)$ for every $0 \leq x \leq \frac{1}{3}$ is a multiplier on $L = [0, 1]$ which is not principal.

Proof. Following Lemma 6, it is suffice to prove that $g_{\frac{1}{3}}$ verify (M_2) , that is, $x \odot (x \rightarrow g_{\frac{1}{3}}(x)) = g_{\frac{1}{3}}(x)$, for every $0 \leq x \leq \frac{1}{3}$.

For $0 \leq x \leq \frac{1}{3}$ we have $\frac{1}{3} \rightarrow x = \max\{\frac{1}{2} - \frac{1}{3}, x\} = \max\{\frac{1}{6}, x\} = \frac{1}{6}$ if $0 \leq x \leq \frac{1}{6}$ and $\frac{1}{3} \rightarrow x = x$ if $\frac{1}{6} < x \leq \frac{1}{3}$, so $g_{\frac{1}{3}}(x) = \frac{1}{3} \odot \frac{1}{6} = 0$ for $0 \leq x \leq \frac{1}{6}$ and $g_{\frac{1}{3}}(x) = \frac{1}{3} \odot x$ for $\frac{1}{6} < x \leq \frac{1}{3}$.

Since for $x > \frac{1}{6}, \frac{1}{3} + x > \frac{1}{3} + \frac{1}{6} = \frac{1}{2}$ we deduce that $g_{\frac{1}{3}}(x) = \frac{1}{3} \odot x = \frac{1}{3} \wedge x = x$, so $g_{\frac{1}{3}}(x) = 0$, for $0 \leq x \leq \frac{1}{6}$ and $g_{\frac{1}{3}}(x) = x$, for $\frac{1}{6} < x \leq \frac{1}{3}$.

Then $x \rightarrow g_{\frac{1}{3}}(x) = x \rightarrow 0 = \max\{\frac{1}{2} - x, 0\} = \frac{1}{2} - x$ for $0 < x \leq \frac{1}{6}$ and $x \rightarrow g_{\frac{1}{3}}(x) = x \rightarrow x = 1$, for $\frac{1}{6} < x \leq \frac{1}{3}$. For $x = 0, 0 \rightarrow g_{\frac{1}{3}}(0) = 1$.

So $x \odot (x \rightarrow g_{\frac{1}{3}}(x)) = x \odot (\frac{1}{2} - x) = 0$, for $0 \leq x \leq \frac{1}{6}$ and $x \odot (x \rightarrow g_{\frac{1}{3}}(x)) = x \odot 1 = x$, for $\frac{1}{6} < x \leq \frac{1}{3}$. For $x = 0, 0 \odot (0 \rightarrow g_{\frac{1}{3}}(0)) = 0$.

We deduce that $x \odot (x \rightarrow g_{\frac{1}{3}}(x)) = g_{\frac{1}{3}}(x)$, for every $0 \leq x \leq \frac{1}{3}$, that is $g_{\frac{1}{3}}$ verify (M_2) , hence $g_{\frac{1}{3}}$ is a multiplier on $L = [0, 1]$. It is easy to prove that $B(L) = \{0, 1\}$, so if suppose by contrary that $g_{\frac{1}{3}}$ is principal, then $g_{\frac{1}{3}} = p_0$ or $g_{\frac{1}{3}} = p_1$ (with $p_0, p_1: [0, \frac{1}{3}] \rightarrow [0, 1]$). Since $g_{\frac{1}{3}}(\frac{1}{3}) = \frac{1}{3} \odot (\frac{1}{3} \rightarrow \frac{1}{3}) = \frac{1}{3} \odot 1 = \frac{1}{3}$ and $p_0(\frac{1}{3}) = 0$ it follows that $g_{\frac{1}{3}} \neq p_0$.

Also, $g_{\frac{1}{3}}(\frac{1}{6}) = \frac{1}{3} \odot (\frac{1}{3} \rightarrow \frac{1}{6}) = \frac{1}{3} \odot 0 = 0$ and $p_1(\frac{1}{6}) = \frac{1}{6}$, so $g_{\frac{1}{3}} \neq p_1$. □

For $J \in \mathcal{I}(L)$, let $M(J, L) = \{p: J \rightarrow L \mid p \text{ is a multiplier on } L\}$, $M(L) = \bigcup\{M(J, L): J \in \mathcal{I}(L)\}$ and $\overline{M}(L) = \{p: L \rightarrow L \mid p \text{ is a multiplier on } L\}$.

Proposition 8 If $J_1, J_2 \in \mathcal{I}(L)$ and $p_i \in M(J_i, L), i = 1, 2$, then

$$(c_{19}) p_1(t) \odot [t \rightarrow p_2(t)] = p_2(t) \odot [t \rightarrow p_1(t)], \text{ for every } t \in J_1 \cap J_2.$$

Proof. For $t \in J_1 \cap J_2$ we have $p_1(t) \odot [t \rightarrow p_2(t)] \stackrel{(M_2)}{=} t \odot (t \rightarrow p_1(t)) \odot (t \rightarrow p_2(t)) = [t \odot (t \rightarrow p_2(t))] \odot (t \rightarrow p_1(t)) \stackrel{(M_2)}{=} p_2(t) \odot [t \rightarrow p_1(t)]$. □

Definition 4 For $J_1, J_2 \in \mathcal{I}(L)$ and $p_i \in M(J_i, L), i = 1, 2$, we define $p_1 \wedge p_2, p_1 \vee p_2, p_1 \otimes p_2, p_1 \rightsquigarrow p_2: J_1 \cap J_2 \rightarrow L$ by $(p_1 \wedge p_2)(t) = p_1(t) \wedge p_2(t), (p_1 \vee p_2)(t) = p_1(t) \vee p_2(t), (p_1 \otimes p_2)(t) = p_1(t) \odot [t \rightarrow p_2(t)] \stackrel{(c_{19})}{=} p_2(t) \odot [t \rightarrow p_1(t)], (p_1 \rightsquigarrow p_2)(t) = t \odot [p_1(t) \rightarrow p_2(t)]$, for every $t \in J_1 \cap J_2$.

Lemma 9 $p_1 \wedge p_2 \in M(J_1 \cap J_2, L)$.

Proof. It is sufficient to verify only (M_2) (for $(M_1), (M_3)$ and (M_4) see Buşneag & Piciu, 2005).

For any $t \in J_1 \cap J_2$ we have $t \odot [t \rightarrow (p_1 \wedge p_2)(t)] = t \odot [t \rightarrow (p_1(t) \wedge p_2(t))] \stackrel{(c_7)}{=} t \odot [(t \rightarrow p_1(t)) \wedge (t \rightarrow p_2(t))] \stackrel{(c_{16})}{=} [t \odot (t \rightarrow p_1(t))] \wedge [t \odot (t \rightarrow p_2(t))] \stackrel{(M_2)}{=} p_1(t) \wedge p_2(t) = (p_1 \wedge p_2)(t)$. □

Lemma 10 $p_1 \vee p_2 \in M(J_1 \cap J_2, L)$.

Proof. The axioms $(M_1), (M_3)$ and (M_4) are verified as in the case of BL algebras (see Buşneag & Piciu, 2005). To verify (M_2) , let $t \in J_1 \cap J_2$. Then $t \odot [t \rightarrow (p_1 \vee p_2)(t)] = t \odot [t \rightarrow (p_1(t) \vee p_2(t))] \stackrel{(c_{13})}{=} t \odot [(t \rightarrow p_1(t)) \vee (t \rightarrow p_2(t))] \stackrel{(c_6)}{=} [t \odot (t \rightarrow p_1(t))] \vee [t \odot (t \rightarrow p_2(t))] \stackrel{(M_2)}{=} p_1(t) \vee p_2(t) = (p_1 \vee p_2)(t)$. □

Lemma 11 $p_1 \otimes p_2 \in M(J_1 \cap J_2, L)$.

Proof. (M_1) is verified as in the case of BL algebras (see Buşneag, & Piciu, 2005), using (c_{10}) . To prove (M_2) , let $t \in J_1 \cap J_2$ and denote $p = p_1 \otimes p_2$.

To prove the equality $t \odot (t \rightarrow p(t)) = p(t)$ it is sufficient to prove that $p(t) \leq t \odot (t \rightarrow p(t))$. We have $p(t) = p_1(t) \odot (t \rightarrow p_2(t)) = t \odot (t \rightarrow p_1(t)) \odot (t \rightarrow p_2(t))$ and $t \odot (t \rightarrow p(t)) = t \odot [t \rightarrow (p_1(t) \odot (t \rightarrow p_2(t)))] = t \odot [t \rightarrow (t \odot (t \rightarrow p_1(t)) \odot (t \rightarrow p_2(t)))]$. So, to prove that $p(t) \leq t \odot (t \rightarrow p(t))$ it is sufficient to prove that $t \odot (t \rightarrow p_1(t)) \odot (t \rightarrow p_2(t)) \leq t \odot [t \rightarrow (t \odot (t \rightarrow p_1(t)) \odot (t \rightarrow p_2(t)))]$, that is, $\varphi \leq t \rightarrow (t \odot \varphi)$ (with $\varphi \stackrel{\text{not}}{=} (t \rightarrow p_1(t)) \odot (t \rightarrow p_2(t))$), which is true, since $\varphi \rightarrow [t \rightarrow (t \odot \varphi)] \stackrel{(c_4)}{=} (\varphi \odot t) \rightarrow (t \odot \varphi) = 1$. (M_3) and (M_4) are verified as in the case of BL algebras (see Buşneag & Piciu, 2005), using (c_{10}) and (c_{11}) . $\square$

Lemma 12 $p_1 \rightsquigarrow p_2 \in M(J_1 \cap J_2, L)$.

Proof. (M_1) is verified as in the case of BL algebras (see Buşneag & Piciu, 2005) using (c_{10}) . To prove (M_2) , let $t \in J_1 \cap J_2$ and denote $p = p_1 \rightsquigarrow p_2 : J_1 \cap J_2 \rightarrow L$; then $p(t) = t \odot [p_1(t) \rightarrow p_2(t)]$. We have $p_1(t) \rightarrow p_2(t) \leq t \rightarrow [t \odot (p_1(t) \rightarrow p_2(t))]$, hence $t \odot [p_1(t) \rightarrow p_2(t)] \leq t \odot [t \rightarrow (t \odot (p_1(t) \rightarrow p_2(t)))] \Leftrightarrow p(t) \leq t \odot [t \rightarrow p(t)] \Leftrightarrow p(t) = t \odot [t \rightarrow p(t)]$.

(M_3) and (M_4) are verified as in the case of BL algebras (see Buşneag, & Piciu, 2005) using (c_{10}) and (c_{11}) . $\square$

Proposition 13

- (i) For every $p \in M(L)$, $p \otimes \mathbf{1} = \mathbf{1} \otimes p = p$;
- (ii) For every $p_1, p_2, p_3 \in M(L)$, $p_1 \otimes (p_2 \otimes p_3) = (p_1 \otimes p_2) \otimes p_3$ and for every $t \in d(p_1) \cap d(p_2) \cap d(p_3)$, $p_1(t) \leq (p_2 \rightsquigarrow p_3)(t)$ iff $(p_1 \otimes p_2)(t) \leq p_3(t)$;
- (iii) For every $p_1, p_2 \in M(L)$ and $t \in d(p_1) \cap d(p_2)$, $(p_1 \rightsquigarrow p_2)(t) \vee (p_2 \rightsquigarrow p_1)(t) = \mathbf{1}(t)$.

Proof. (i) If $J = \text{dom}(p)$ and $t \in J$, then $(p \otimes \mathbf{1})(t) = p(t) \odot (t \rightarrow \mathbf{1}(t)) = p(t) \odot (t \rightarrow t) = p(t) \odot 1 = p(t)$ and $(\mathbf{1} \otimes p)(t) = t \odot (t \rightarrow p(t)) = p(t)$, that is, $p \otimes \mathbf{1} = \mathbf{1} \otimes p = p$.

(ii) Let $p_i \in M(J_i, L)$ where $J_i \in \mathcal{I}(L)$, $i = 1, 2, 3$. Thus, for $t \in J_1 \cap J_2 \cap J_3$ we have $[p_1 \otimes (p_2 \otimes p_3)](t) = ((p_2 \otimes p_3)(t)) \odot (t \rightarrow p_1(t)) = [p_2(t) \odot (t \rightarrow p_3(t))] \odot (t \rightarrow p_1(t)) = p_2(t) \odot [(t \rightarrow p_3(t)) \odot (t \rightarrow p_1(t))] = p_2(t) \odot [(t \rightarrow p_1(t)) \odot (t \rightarrow p_3(t))] = [p_2(t) \odot (t \rightarrow p_1(t))] \odot (t \rightarrow p_3(t)) = ((p_1 \otimes p_2)(t)) \odot (t \rightarrow p_3(t)) = [(p_1 \otimes p_2) \otimes p_3](t)$, that is the operation $\otimes$ is associative.

For $t \in J_1 \cap J_2 \cap J_3$ we have $p_1(t) \leq (p_2 \rightsquigarrow p_3)(t) \Leftrightarrow p_1(t) \leq t \odot [p_2(t) \rightarrow p_3(t)]$. So, by (c_3) , $p_1(t) \odot [t \rightarrow p_2(t)] \leq t \odot (t \rightarrow p_2(t)) \odot (p_2(t) \rightarrow p_3(t)) \stackrel{(M_2)}{\Leftrightarrow} p_1(t) \odot [t \rightarrow p_2(t)] \leq p_2(t) \odot (p_2(t) \rightarrow p_3(t)) \leq p_3(t) \Leftrightarrow (p_1 \otimes p_2)(t) \leq p_3(t)$, for any $t \in J_1 \cap J_2 \cap J_3$, that is, $p_1 \otimes p_2 \leq p_3$. Conversely, if $(p_1 \otimes p_2)(t) \leq p_3(t)$ we have $p_2(t) \odot [t \rightarrow p_1(t)] \leq p_3(t)$, for any $t \in J_1 \cap J_2 \cap J_3$. Obviously, $t \rightarrow p_1(t) \leq p_2(t) \rightarrow p_3(t) \stackrel{(c_3)}{\Leftrightarrow} t \odot (t \rightarrow p_1(t)) \leq t \odot (p_2(t) \rightarrow p_3(t)) \Leftrightarrow p_1(t) \leq (p_2 \rightsquigarrow p_3)(t)$.

(iii) We have $(p_1 \rightsquigarrow p_2)(t) \vee (p_2 \rightsquigarrow p_1)(t) = [t \odot (p_1(t) \rightarrow p_2(t))] \vee [t \odot (p_2(t) \rightarrow p_1(t))] = t \odot [(p_1(t) \rightarrow p_2(t)) \vee (p_2(t) \rightarrow p_1(t))] = t \odot 1 = t = \mathbf{1}(t)$. $\square$

Corollary 14 $(\overline{M}(L), \wedge, \vee, \otimes, \rightsquigarrow, \mathbf{0}, \mathbf{1})$ is a MTL algebra.

Definition 5 (Esteva & Godo, 2001; Freytes, 2004) A MTL algebra L is called

- (i) an $IMTL$ algebra (involutive algebra) if it satisfies the equation

$$(I) x^{**} = x;$$
- (ii) a $SMTL$ algebra if it satisfies the equation

$$(S) x \wedge x^* = 0;$$
- (iii) a WNM algebra (weak nilpotent minimum) if it satisfies the equation

$$(W) (x \odot y)^* \vee [(x \wedge y) \rightarrow (x \odot y)] = 1;$$
- (iv) a $\Pi S MTL$ algebra if it is a $SMTL$ algebra satisfying the equation

$$(\Pi) [z^{**} \odot ((x \odot z) \rightarrow (y \odot z))] \rightarrow (x \rightarrow y) = 1.$$

Theorem 15

- (i) If L is a BL algebra, then for every $f_1, f_2 \in M(L)$, $(f_1 \otimes (f_1 \rightsquigarrow f_2))(x) = (f_1 \wedge f_2)(x)$, for every $x \in d(f_1) \cap d(f_2)$;
- (ii) If L is an $IMTL$ algebra, then $f^{**} = f$, for every $f \in M(L)$;
- (iii) If L is a $SMTL$ algebra, then for every $f \in \overline{M}(L)$, $f \wedge f^* = \mathbf{0}$;
- (iv) If L is a WNM algebra, then for every $f_1, f_2 \in \overline{M}(L)$, $(f_1 \otimes f_2)^* \vee ((f_1 \wedge f_2) \rightsquigarrow (f_1 \otimes f_2)) = \mathbf{1}$;
- (v) If L is a $\Pi SMTL$ algebra, then for every $f, g, h \in \overline{M}(L)$, $[h^{**} \otimes ((f \otimes h) \rightsquigarrow (g \otimes h))] \rightsquigarrow (f \rightsquigarrow g) = \mathbf{1}$.

Proof. (i) Suppose L is a BL algebra (see Remark 1). Let $f_1, f_2 \in M(L)$, $f_1: J_1 \rightarrow L$, $f_2: J_2 \rightarrow L$, with $J_1, J_2 \in \mathcal{I}(L)$. For every $x \in J_1 \cap J_2$ we have $(f_1 \otimes (f_1 \rightsquigarrow f_2))(x) = (f_1 \wedge f_2)(x) \Leftrightarrow (f_1 \rightsquigarrow f_2)(x) \odot [x \rightarrow f_1(x)] = f_1(x) \wedge f_2(x) \Leftrightarrow x \odot [f_1(x) \rightarrow f_2(x)] \odot [x \rightarrow f_1(x)] = f_1(x) \wedge f_2(x) \Leftrightarrow (x \odot [x \rightarrow f_1(x)]) \odot [f_1(x) \rightarrow f_2(x)] = f_1(x) \wedge f_2(x) \stackrel{(M_2)}{\Leftrightarrow} f_1(x) \odot [f_1(x) \rightarrow f_2(x)] = f_1(x) \wedge f_2(x)$, which is true because L is supposed a BL algebra.

(ii) Suppose L is an $IMTL$ algebra. For $f \in M(L)$, $f: J \rightarrow L$ and $x \in J$, we have $f^{**} = (f \rightsquigarrow \mathbf{0}) \rightsquigarrow \mathbf{0}$ and $f^{**}(x) = x \odot [x \odot f^*(x)]^* \stackrel{(c_4)}{=} x \odot [x \rightarrow (f(x))^{**}] \stackrel{(I)}{=} x \odot [x \rightarrow f(x)] \stackrel{(M_2)}{=} f(x)$, hence $f^{**} = f$.

(iii) Suppose L is a $SMTL$ algebra. If $f \in \overline{M}(L)$, $f: L \rightarrow L$, then the equation $f \wedge f^* = \mathbf{0}$ is equivalent with $f \wedge (f \rightsquigarrow \mathbf{0}) = \mathbf{0} \Leftrightarrow f(x) \wedge [x \odot (f(x))^*] = 0$, for every $x \in L$, which is clearly (since $f(x) \wedge [x \odot (f(x))^*] \leq f(x) \wedge (f(x))^* = 0$), hence $f \wedge f^* = \mathbf{0}$.

(iv) Suppose L is a WNM algebra. Let $f_1, f_2 \in \overline{M}(L)$, $f_1, f_2: L \rightarrow L$ and $x \in L$. We denote $a = f_1(x)$, $b = f_2(x)$. We have $((f_1 \otimes f_2)^* \vee ((f_1 \wedge f_2) \rightsquigarrow (f_1 \otimes f_2)))(x) = ((f_1 \otimes f_2)^*(x)) \vee (x \odot ((f_1 \wedge f_2)(x) \rightarrow (f_1 \otimes f_2)(x))) = (x \odot (a \odot (x \rightarrow b))^*) \vee (x \odot ((a \wedge b) \rightarrow (a \odot (x \rightarrow b)))) \stackrel{(c_6)}{=} x \odot ((a \odot (x \rightarrow b))^* \vee ((a \wedge b) \rightarrow (a \odot (x \rightarrow b))))$.

Since $b \leq x \rightarrow b$ we deduce that $a \wedge b \leq a \wedge (x \rightarrow b)$, hence (using (c_3)) $(a \wedge (x \rightarrow b)) \rightarrow (a \odot (x \rightarrow b)) \leq (a \wedge b) \rightarrow (a \odot (x \rightarrow b))$.

Since L is supposed a WNM -algebra we obtain $1 = (a \odot (x \rightarrow b))^* \vee ((a \wedge (x \rightarrow b)) \rightarrow (a \odot (x \rightarrow b))) \leq (a \odot (x \rightarrow b))^* \vee ((a \wedge b) \rightarrow (a \odot (x \rightarrow b)))$, hence $(a \odot (x \rightarrow b))^* \vee ((a \wedge b) \rightarrow (a \odot (x \rightarrow b))) = 1$. Then $((f_1 \otimes f_2)^* \vee ((f_1 \wedge f_2) \rightsquigarrow (f_1 \otimes f_2)))(x) = x \odot 1 = x = \mathbf{1}(x) \Leftrightarrow (f_1 \otimes f_2)^* \vee ((f_1 \wedge f_2) \rightsquigarrow (f_1 \otimes f_2)) = \mathbf{1}$.

(v) Suppose now L is a $\Pi SMTL$ algebra. From the condition $x \wedge x^* = 0$ ($x \in L$), we deduce that $x^* \vee x^{**} \stackrel{(c_{15})}{=} (x \wedge x^*)^* = 0^* = 1$, that is, $x^* \in B(L)$. For $f, g, h: L \rightarrow L$, and $x \in L$ we denote $a_1 = f(x)$, $a_2 = g(x)$ and $a_3 = h(x)$. Then $h^{**}(x) = x \odot (x \rightarrow a_3^{**}) \stackrel{(c_9)}{=} x \wedge a_3^{**} \stackrel{(c_8)}{=} x \odot a_3^{**}$, $[h^{**} \otimes ((f \otimes h) \rightsquigarrow (g \otimes h))](x) = [x \rightarrow h^{**}(x)] \odot [x \odot ((f \otimes h)(x) \rightarrow (g \otimes h)(x))] = [x \rightarrow (x \odot a_3^{**})] \odot [x \odot (((x \rightarrow a_1) \odot a_3) \rightarrow ((x \rightarrow a_2) \odot a_3))] \stackrel{(c_1)}{=} [x \odot (x \rightarrow (x \odot a_3^{**}))] \odot [((x \rightarrow a_1) \odot a_3) \rightarrow ((x \rightarrow a_2) \odot a_3)] \leq (x \odot a_3^{**}) \odot [((x \rightarrow a_1) \odot a_3) \rightarrow ((x \rightarrow a_2) \odot a_3)] \stackrel{(II)}{=} x \odot [a_3^{**} \odot (((x \rightarrow a_1) \odot a_3) \rightarrow ((x \rightarrow a_2) \odot a_3))] \leq x \odot [(x \rightarrow a_1) \rightarrow (x \rightarrow a_2)] = x \odot [(x \odot (x \rightarrow a_1)) \rightarrow a_2] \stackrel{(M_2)}{=} x \odot (a_1 \rightarrow a_2) = (f \rightsquigarrow g)(x)$, hence $[h^{**} \otimes ((f \otimes h) \rightsquigarrow (g \otimes h))] \rightsquigarrow (f \rightsquigarrow g) = \mathbf{1}$. $\square$

Corollary 16 If L is a BL algebra (resp. an $IMTL$ algebra, a $SMTL$ algebra, a WNM algebra, a $\Pi SMTL$ algebra) then $\overline{M}(L)$ is a BL algebra (resp. an $IMTL$ algebra, a $SMTL$ algebra, a WNM algebra, a $\Pi SMTL$ algebra).

Using the rules (c_6) , (c_{10}) and (c_{11}) we obtain:

Lemma 17 $v_L: B(L) \rightarrow \overline{M}(L)$, $v_L(a) = \overline{p}_a$ for every $a \in B(L)$, is a monomorphism of MTL algebras.

Definition 6 A subset $J \subseteq L$ is called regular if for every $x, y \in L$ such that $x \wedge f = y \wedge f$ for every $f \in B(L) \cap J$, then $x = y$.

Denote $\text{Reg}(L)$ the set of all regular subset of L .

Example 3 We give an example of non-trivial regular subset in a MTL algebra. Consider $L = \{0, a, b, c, d, 1\}$ the MTL -algebra from Example 1. We have that $B(L) = \{0, a, d, 1\}$ and if consider $J = \{0, b, d, a\}$, then $J \in \mathcal{I}(L)$. It is easy to prove that for any $x, y \in L$ with $x \neq y$, there is $f \in J \cap B(L) = \{0, a, d\}$ such that $x \wedge f \neq y \wedge f$, that is, $J \in \text{Reg}(L)$.

Remark 5 The condition $J \in \text{Reg}(L)$ is equivalent with: if $x, y \in L$ and $p_{x|J \cap B(L)} = p_{y|J \cap B(L)}$, then $x = y$.

Lemma 18 If $J_1, J_2 \in \mathcal{I}(L) \cap \text{Reg}(L)$, then $J_1 \cap J_2 \in \mathcal{I}(L) \cap \text{Reg}(L)$.

Denote $M_{reg}(L) = \{p \in M(L): d(p) \in I(L) \cap Reg(L)\}$.

Remark 6 By Lemmas 9-12 and 18 we deduce that if $p_1, p_2 \in M_{reg}(L)$, then $p_1 \otimes p_2, p_1 \rightsquigarrow p_2 \in M_{reg}(L)$.

Proposition 19 Let $p: J \rightarrow L$ be a multiplier on L with $J \in I(L) \cap Reg(L)$. Then $(p \vee p^*)(x) = x$, for every $x \in J$.

Proof. Let $a \in B(L) \cap J$ and $x \in J$. Then

$$\begin{aligned}
 a \wedge [p \vee p^*](x) &= a \wedge [p(x) \vee (x \odot (p(x))^*)] = [a \wedge p(x)] \vee [a \wedge (x \odot (p(x))^*)] \\
 &= [x \odot p(a)] \vee [x \odot a \odot (p(x))^*] \stackrel{c_{10}}{=} [x \odot p(a)] \vee [x \odot a \odot (a \odot p(x))^*] \\
 &\stackrel{M_4}{=} [x \odot p(a)] \vee [x \odot a \odot (x \odot p(a))^*] = [x \odot p(a)] \vee [x \odot a \odot (x \wedge p(a))^*] \\
 &\stackrel{c_{15}}{=} [x \odot p(a)] \vee [x \odot a \odot (x^* \vee (p(a))^*)] \stackrel{c_6}{=} [x \odot p(a)] \vee [a \odot ((x \odot x^*) \vee (x \odot (p(a))^*))] \\
 &= [x \odot p(a)] \vee [a \odot (0 \vee (x \odot (p(a))^*))] = [x \odot p(a)] \vee [a \odot x \odot (p(a))^*] \\
 &= [x \odot p(a)] \vee [x \odot (a \odot (p(a))^*)] \stackrel{c_6}{=} x \odot [p(a) \vee (a \odot (p(a))^*)] \\
 &= x \odot [p(a) \vee (a \wedge (p(a))^*)] = x \odot [(p(a) \vee a) \wedge (p(a) \vee (p(a))^*)] \\
 &= x \odot (a \wedge 1) = x \odot a = x \wedge a,
 \end{aligned}$$

so $(p \vee p^*)(x) = x$, since $J \in Reg(L)$. $\square$

Definition 7 Let two multipliers p_1, p_2 on L . We say that p_2 extends p_1 if $d(p_1) \subseteq d(p_2)$ and $p_{2|d(p_1)} = p_1$; if p_2 extends p_1 , we write $p_1 \sqsubseteq p_2$. If we can not be extended a multiplier p to a strictly larger domain, we called p maximal.

Lemma 20

(i) If $p_1, p_2 \in M(L)$, $p \in M_{reg}(L)$ and $p \sqsubseteq p_1, p \sqsubseteq p_2$, then p_1 and p_2 coincide on the $d(p_1) \cap d(p_2)$;

(ii) any $p \in M_r(L)$ can be extended to a maximal multiplier. For any principal multiplier $p_a, a \in B(L), d(p_a) \in I(L) \cap Reg(L)$ there is a uniquely total multiplier $\overline{p_a}$ such that $p_a \sqsubseteq \overline{p_a}$ and for any non-principal multiplier p there is a maximal non-principal multiplier r such that $p \sqsubseteq r$.

On $M_{reg}(L)$ we consider the relation ρ_L defined by $(p_1, p_2) \in \rho_L$ iff $p_{1|d(p_1) \cap d(p_2)} = p_{2|d(p_1) \cap d(p_2)}$.

Lemma 21 ρ_L is an equivalence relation on $M_{reg}(L)$ compatible with $\wedge, \vee, \otimes$ and $\rightsquigarrow$.

Proof. Obviously, ρ_L is an equivalence relation on $M_{reg}(L)$ compatible with $\wedge$ and $\vee$.

For the compatibility of ρ_L with $\otimes$ and $\rightsquigarrow$ on $M_{reg}(L)$, let $(p_1, p_2), (r_1, r_2) \in \rho_L$.

Let $t \in d(p_1) \cap d(p_2) \cap d(r_1) \cap d(r_2)$. We have $p_1(t) = p_2(t)$ and $r_1(t) = r_2(t)$, so

$$\begin{aligned}
 (p_1 \otimes r_1)(t) &= p_1(t) \odot (t \rightarrow r_1(t)) = p_2(t) \odot (t \rightarrow r_2(t)) = (p_2 \otimes r_2)(t), \\
 (p_1 \rightsquigarrow r_1)(t) &= t \odot [p_1(t) \rightarrow r_1(t)] = t \odot [p_2(t) \rightarrow r_2(t)] = (p_2 \rightsquigarrow r_2)(t),
 \end{aligned}$$

that is, $(p_1 \otimes r_1, p_2 \otimes r_2), (p_1 \rightsquigarrow r_1, p_2 \rightsquigarrow r_2) \in \rho_L$. $\square$

For $p \in M_{reg}(L)$ with $J = d(p) \in I(L) \cap Reg(L)$, we denote by $[p, J]$ the congruence class of p modulo ρ_L and $L'' = M_{reg}(L)/\rho_L$.

On L'' we define the order relation $[p_1, J_1] \leq [p_2, J_2]$ iff $p_1(x) \leq p_2(x)$, for every $x \in J_1 \cap J_2$.

It is a routine to prove the following result:

Lemma 22 $(L'', \leq)$ is a bounded lattice, where for $[p_1, J_1], [p_2, J_2] \in L'', [p_1, J_1] \wedge [p_2, J_2] = [p_1 \wedge p_2, J_1 \cap J_2]$ and $[p_1, J_1] \vee [p_2, J_2] = [p_1 \vee p_2, J_1 \cap J_2], \mathbf{0} = [\mathbf{0}, L], \mathbf{1} = [\mathbf{1}, L]$.

For $[p_1, J_1], [p_2, J_2] \in L''$, we define $[p_1, J_1] \otimes [p_2, J_2] = [p_1 \otimes p_2, J_1 \cap J_2]$ and $[p_1, J_1] \rightsquigarrow [p_2, J_2] = [p_1 \rightsquigarrow p_2, J_1 \cap J_2]$ (where $[p_1 \otimes p_2]$ and $[p_1 \rightsquigarrow p_2]$ are defined in Definition 4).

Proposition 23 $(L'', \wedge, \vee, \otimes, \rightsquigarrow, \mathbf{0}, \mathbf{1})$ is a MTL-algebra.

Proof. We verify the axioms of MTL-algebras.

(LR₁) Follows from Lemma 22;

(LR₂) Follows from Proposition 13, (i), (ii);

(LR₃) Follows from Proposition 13, (ii);

The prelinarity equation (c₁₂) follows from Proposition 13, (iii). $\square$

Remark 7 From Theorem 2, Propositions 19 and 23 we deduce that L'' is a Boolean algebra.

Remark 8 If consider $\mathcal{F} = I(L) \cap \text{Reg}(L)$ and the partially ordered systems $\{\delta_{I,J}\}_{I,J \in \mathcal{F}, I \subseteq J}$ (for $I, J \in \mathcal{F}, I \subseteq J, \delta_{I,J}: M(J, L) \rightarrow M(I, L)$ is defined by $\delta_{I,J}(f) = f|_I$), then $L'' = \lim_{I \in \mathcal{F}} M(I, L)$.

Lemma 24 If consider $\bar{v}_L: B(L) \rightarrow L''$ defined by $\bar{v}_L(e) = [\bar{p}_e, L]$ for any $e \in B(L)$, then:

(i) $\bar{v}_L$ is a monomorphism of Boolean algebras;

(ii) $\bar{v}_L(B(L)) \in \text{Reg}(L'')$.

Proof. (i) See Lemma 17.

(ii) To prove $\bar{v}_L(B(L)) \in \text{Reg}(L'')$, we suppose by contrary that there exist $p_1, p_2 \in M_{\text{reg}}(L)$ such that $[p_1, d(p_1)] \neq [p_2, d(p_2)]$ (hence we have $a \in d(p_1) \cap d(p_2)$ such that $p_1(a) \neq p_2(a)$) and $[p_1, d(p_1)] \wedge [\bar{p}_e, L] = [p_2, d(p_2)] \wedge [\bar{p}_e, L]$ for any $[\bar{p}_e, L] \in \bar{v}_L(B(L)) \cap B(L'') \Rightarrow p_1(x) \wedge e \wedge x = p_2(x) \wedge e \wedge x$ for any $x \in d(p_1) \cap d(p_2)$ and any $e \in B(L)$. For $e = 1$ and $x = a$ we deduce that $p_1(a) \wedge a = p_2(a) \wedge a \Leftrightarrow p_1(a) = p_2(a)$, a contradiction. $\square$

Remark 9 Following Lemma 20 we can identify $[\bar{p}_e, L]$ with $\bar{p}_e$, for every $e \in B(L)$. So, the boolean elements can be identified with the elements of $\{\bar{p}_e: e \in B(L)\}$.

Following the above consideration we deduce, as in the case of BL -algebras (see Buşneag & Piciu, 2005), that:

Lemma 25 If $[p, d(p)] \in L''$ (with $p \in M_{\text{reg}}(L)$ and $J = d(p) \in I(L) \cap \text{Reg}(L)$), then $J \cap B(L) \subseteq \{e \in B(L): \bar{p}_e \wedge [p, d(p)] \in B(L)\}$.

4. Maximal MTL Algebra of Quotients

In this section by L we denote a MTL -algebra.

Definition 8 A MTL algebra G is called MTL algebra of fractions of L if:

(Fr₁) $B(L)$ is a MTL subalgebra of G ;

(Fr₂) For every $f, g, h \in G, f \neq g$, there is $e \in B(L)$ such that $e \wedge f \neq e \wedge g$ and $e \wedge h \in B(L)$.

We write $L \sqsubseteq G$ if G is a MTL algebra of fractions of L .

Definition 9 $Q(L)$ is the maximal MTL algebra of quotients of L if $L \sqsubseteq Q(L)$ and for any MTL algebra G with $L \sqsubseteq G$ there is a injective morphism of MTL algebras $j: G \rightarrow Q(L)$.

Proposition 26 If L is a MTL -algebra and $L \sqsubseteq G$, then G is a Boolean algebra.

Proof. Indeed, if suppose that G is not a Boolean algebra, by Theorem 2, there is $f \in G$ such that $f \vee f^* \neq 1$. Since $L \sqsubseteq G$, then there is a boolean element g such that $g \wedge f$ is boolean and $g \wedge (f \vee f^*) \neq g$. Since $g \wedge f \in B(L)$, then $(g \wedge f) \vee (g \wedge f)^* = 1 \Rightarrow (g \wedge f) \vee (g^* \vee f^*) = 1 \Rightarrow [(g \wedge f) \vee g^*] \vee f^* = 1 \Rightarrow [(g \vee g^*) \wedge (f \vee g^*)] \vee f^* = 1 \Rightarrow [1 \wedge (f \vee g^*)] \vee f^* = 1 \Rightarrow (f \vee f^*) \vee g^* = 1$. By the unicity of the complement of g , we deduce that $f \vee f^* = g$. Then from $g \wedge (f \vee f^*) \neq g$ we obtain $g \wedge g \neq g \Rightarrow g \neq g$, a contradiction. Hence G is a Boolean algebra. $\square$

Corollary 27 $Q(L)$ is a Boolean algebra.

As in Buşneag and Piciu (2005), we have:

Remark 10 If L is a Boolean algebra, obviously, $B(L) = L$. By Proposition 26, $Q(L)$ is a Boolean algebra; the axioms M_1, M_2, M_3 are equivalent with M_4 and $Q(L)$ is just Dedekind-MacNeille completion of L (Schmid, 1980).

Lemma 28 Let $L \sqsubseteq G$; then for every $f, g \in G, f \neq g, h_1, \dots, h_n \in G$, there exists a boolean element e such that $e \wedge f \neq e \wedge g$ and $e \wedge h_i \in B(L)$ for $i = 1, 2, \dots, n$ ($n \geq 2$).

Lemma 29 Let $L \sqsubseteq G$ and $g \in G$. Then $I_g = \{e \in B(L): e \wedge g \in B(L)\} \in I(B(L)) \cap M_{\text{reg}}(L)$.

Theorem 30 L'' (defined in Section 3) is the maximal (boolean) MTL algebra $Q(L)$ of quotients of L .

Proof. From Lemma 24, (i), $B(L)$ is a MTL subalgebra of L'' . Consider $[f_1, d(f_1)], [f_2, d(f_2)], [f_3, d(f_3)] \in L''$, $f_1, f_2, f_3 \in M_{\text{reg}}(L)$ such that $[f_2, d(f_2)] \neq [f_3, d(f_3)]$. Then we have $x' \in d(f_2) \cap d(f_3)$ with $f_2(x') \neq f_3(x')$.

Consider $J = d(f_1) \in I(L) \cap \text{Reg}(L)$ and $J_{[f_1, d(f_1)]} = \{a \in B(L): \bar{p}_a \wedge [f_1, d(f_1)] \in B(L)\}$. From Lemma 25, $J \cap B(L) \subseteq I_{[f_1, d(f_1)]}$. If suppose that for any $a \in J \cap B(L)$, $\bar{p}_a \wedge [f_2, d(f_2)] = \bar{p}_a \wedge [f_3, d(f_3)]$, then $[\bar{p}_a \wedge f_2, d(f_2)] =$

$[\overline{p_a} \wedge f_3, d(f_3)]$, so for any $x \in d(f_2) \cap d(f_3)$ we have $(\overline{p_a} \wedge f_2)(x) = (\overline{p_a} \wedge f_3)(x)$ i.e. $a \wedge f_2(x) = a \wedge f_3(x)$. Because $J \in \text{Re } g(L)$, $f_2(x) = f_3(x)$ for any $x \in d(f_2) \cap d(f_3)$ so $[f_2, d(f_2)] = [f_3, d(f_3)]$, a contradiction.

If $[f_2, d(f_2)] \neq [f_3, d(f_3)]$, then there is $a \in J \cap B(L)$, such that $\overline{p_a} \wedge [f_2, d(f_2)] \neq \overline{p_a} \wedge [f_3, d(f_3)]$.

Since by Lemma 25, $J \cap B(L) \subseteq J_{[f_1, d(f_1)]}$ for this $a \in J \cap B(L)$ we have $\overline{p_a} \wedge [f_1, d(f_1)] \in B(L)$.

Now, consider G a MTL algebra such that $L \sqsubseteq G$; obviously, $B(L) \subseteq B(G)$

$$\begin{array}{ccc} L & \sqsubseteq & G \\ & \swarrow j & \\ & L'' & \end{array}$$

By Lemma 29, For $a' \in G$, $J_{a'} = \{e \in B(L) : e \wedge a' \in B(L)\} \in I(B(L)) \cap \text{Reg}(L)$.

$p_{a'} : J_{a'} \rightarrow L$, $p_{a'}(x) = x \wedge a'$ is a multiplier. Indeed, (M_1) and (M_2) are verified, because if $e \in B(L)$ and $x \in J_{a'}$, then $p_{a'}(e \odot x) = (e \odot x) \wedge a' = (e \wedge x) \wedge a' = e \wedge (x \wedge a') = e \odot (x \wedge a') = e \odot p_{a'}(x)$, and $x \odot (x \rightarrow p_{a'}(x)) = x \odot [x \rightarrow (x \wedge a')] \stackrel{(c_9)}{=} x \wedge (x \wedge a') = x \wedge a' = p_{a'}(x)$. To verify (M_3) , let $e \in J_{a'} \cap B(L) = J_{a'}$. Thus, $p_{a'}(e) = e \wedge a' \in B(L)$ (since $e \in J_{a'}$). The condition (M_4) is obviously verified, hence $[p_{a'}, J_{a'}] \in L''$.

We define $j : G \rightarrow L''$, by $j(a') = [p_{a'}, J_{a'}]$, for every $a' \in G$. Obviously, $j(0) = \mathbf{0}$. For $a', b' \in G$ and $x \in J_{a'} \cap J_{b'}$, $(j(a') \otimes j(b'))(x) = (a' \wedge x) \odot [x \rightarrow (b' \wedge x)] = (a' \odot x) \odot [x \rightarrow (b' \wedge x)] = a' \odot [x \odot (x \rightarrow (b' \wedge x))] = a' \odot [x \wedge (b' \wedge x)] = a' \odot (b' \wedge x) = a' \odot (b' \odot x) = (a' \odot b') \odot x = (a' \odot b') \wedge x = j(a' \odot b')(x)$, hence $j(a') \otimes j(b') = j(a' \odot b')$ and $(j(a') \rightsquigarrow j(b'))(x) = x \odot [j(a')(x) \rightarrow j(b')(x)] = x \odot [(a' \wedge x) \rightarrow (b' \wedge x)] = x \odot [(x \odot a') \rightarrow (x \odot b')] \stackrel{(c_{10})}{=} x \odot (a' \rightarrow b') = x \wedge (a' \rightarrow b') = j(a' \rightarrow b')(x)$, hence $j(a') \rightsquigarrow j(b') = j(a' \rightarrow b')$.

Now, let $a', b' \in G$ such that $j(a') = j(b')$. It follows that $[p_{a'}, J_{a'}] = [p_{b'}, J_{b'}]$, so $p_{a'}(x) = p_{b'}(x)$ for any $x \in J_{a'} \cap J_{b'}$. So $a' \wedge x = b' \wedge x$ for any $x \in J_{a'} \cap J_{b'}$. By Lemma 28, if $a' \neq b'$, since $L \sqsubseteq G$, there is a boolean element e such that $e \wedge a', e \wedge b' \in B(L)$ and $e \wedge a' \neq e \wedge b'$ which is contradictory (since $e \wedge a', e \wedge b' \in B(L)$ implies $e \in J_{a'} \cap J_{b'}$). $\square$

Proposition 31 Let L be a MTL-algebra. The following are equivalent:

- (i) Every maximal multiplier on L has domain L ;
- (ii) For every multiplier $p \in M(J, L)$ there is $e \in B(L)$ such that $p = p_e$, (i.e., $p(x) = e \wedge x$ for any $x \in J$);
- (iii) $Q(L) \approx B(L)$.

Proof. (i) $\Rightarrow$ (ii) Assume (i) and for $p \in M(J, L)$ let p' its the maximal extension (by Lemma 20). By (i), we have $p' : L \rightarrow L$. Put $e = p'(1) \in B(L)$ (by M_3), then for every $x \in J$, $p(x) = p(x) \wedge 1 \stackrel{M_4}{=} x \wedge p(1) = x \wedge e = p_e(x)$, that is $p = p_e$.

(ii) $\Rightarrow$ (iii) Follow from Lemma 24.

(iii) $\Rightarrow$ (i) Follow from Lemma 20 and Lemma 24. $\square$

Remark 11

1) If L is a MTL algebra with $B(L) = L_2$ and $L \sqsubseteq G$ then $G = \{0, 1\}$, hence $Q(L) = L'' \approx L_2$. Indeed, if $a_1, a_2, a_3 \in G, a_1 \neq a_2$, then there exists $e \in B(L)$ (by (Fr_2)) such that $e \wedge a_1 \neq e \wedge a_2$ (hence $e \neq 0$) and $e \wedge a_3 \in B(L)$. Clearly, $e = 1$, hence $a_3 \in B(L)$, that is, $G = B(L)$.

2) More general, if L is a MTL-algebra such that $B(L)$ is finite and $L \sqsubseteq G$ then $G = B(L)$, hence $Q(L) = B(L)$. Indeed, since $L \sqsubseteq G$, we have $B(L) \subseteq G$. Let $a \in G$. Then there is $e \in B(L)$ such that $e \wedge a \in B(L)$. $Q(L)$ is finite, so, there is a largest element $e_a \in Q(L)$ with $e_a \wedge a \in B(L)$. Suppose $e_a \vee a \neq e_a$. Then there is $e \in B(L)$ such that $e \wedge (e_a \vee a) \neq e \wedge e_a$ and $e \wedge a \in B(L)$. Because $e \wedge a \in B(L)$ we deduce $e \leq e_a$ so $e = e \wedge (e_a \vee a) \neq e \wedge e_a = e$, a contradiction. Hence $e_a \vee a = e_a$, so $a \leq e_a$, consequently $a = a \wedge e_a \in Q(L)$, that is $G \subseteq Q(L)$. Then $G = Q(L)$, hence $Q(L) \approx B(L)$.

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